# CS3191 Section 4 Large Games

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There are three problems which have to be solved to write such a program which we will discuss in some detail. Finally we will have a look at how Chess-playing programs developed, since Chess is the game for which the most effort has been made when it comes to writing programs.

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Board representation and move generation. Clearly we have to think about how the board (and the pieces) are represented internally, and how the moves are to be generated. Typically, once this has been solved it can be left alone.

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Board representation and move generation.

Alpha-beta search. Despite the fact that we cannot hope to employ the minimax algorithm with alpha-beta pruning, this technique still plays a vital role in game-playing programs. There are some variants that might be implemented, and typically some effort is spent on cataloguing search results.

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Alpha-beta search.

Evaluation function. Since alpha-beta search cannot be carried out until a leaf is reached, the search stops instead at a pre-defined depth. To obtain a value for a position at this depth, a function has to be created which assigns one based entirely on the state of the board at the time. This is known as the 'evaluation function'.

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The faster the program, the higher the depth to which it can carry out alpha-beta search (before it has to 'guess' a value for a position), and the better it will play. Hence speed is of the essence when writing such programs, and is a concern for all the components mentioned above.

# Representing the board and related issues

# Representing the board-array

In order to illustrate our thoughts, we often use Chess as an example. However, there's no need to be familiar with the game beyond the rudiments.

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To generate moves: Pick piece, generate possible target fields, then:

- check target field not occupied by own piece;
- if piece is a rook, bishop, pawn or queen, check whether the way to target is empty;
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#### Need:

- loop over all fields (to pick piece);
- loop over all possible target positions;
- loop to check for obstructions along the way.

Complicated, not fast.

# **Board representation**-0x88

Faster: Assign a number to each square on the board given by one byte, four high bits: row; four low bits: column.

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|           | а  | b  | С   | d  | е   | f   | g   | h   |   |
|-----------|--|--|---|--|---|---|---|---|---|
|           | 0000   | 0001   | 0010  | 0011   | 0100  | 0101  | 0110  | 0111  | low   |
| 0111      | 112  | 113  | 114   | 115  | 116   | 117   | 118   | 119   |   |
| 0110      | 96   | 97   | 98  | 99   | 100   | 101   | 102   | 103   |   |
| 0101      | 80   | 81   | 82  | 83   | 84  | 85  | 86  | 87  |   |
| 0100      | 64   | 65   | 66  | 67   | 68  | 69  | 70  | 71  |   |
| 0011      | 48   | 49   | 50  | 51   | 52  | 53  | 54  | 55  |   |
| 0010      | 32   | 33   | 34  | 35   | 36  | 37  | 38  | 39  |   |
| 0001      | 16   | 17   | 18  | 19   | 20  | 21  | 22  | 23  |   |
| 0000      | 0  | 1  | 2   | 3  | 4   | 5   | 6   | 7   |   |
| high bits |  |  |   |  |   |   |   |   | •   |
|           | 0110<br>0101<br>0100<br>0011<br>0010<br>0001<br>0000 | 0000       0111     112       0110     96       0101     80       0100     64       0011     48       0010     32       0001     16       0000     0 | 0000     0001       0111     112     113       0110     96     97       0101     80     81       0100     64     65       0011     48     49       0010     32     33       0001     16     17       0000     0     1 | 0000     0001     0010       0111     112     113     114       0110     96     97     98       0101     80     81     82       0100     64     65     66       0011     48     49     50       0010     32     33     34       0001     16     17     18       0000     0     1     2 | 0000         0001         0010         0011           0111         112         113         114         115           0110         96         97         98         99           0101         80         81         82         83           0100         64         65         66         67           0011         48         49         50         51           0010         32         33         34         35           0001         16         17         18         19           0000         0         1         2         3 | 0000         0001         0010         0011         0100           0111         112         113         114         115         116           0110         96         97         98         99         100           0101         80         81         82         83         84           0100         64         65         66         67         68           0011         48         49         50         51         52           0010         32         33         34         35         36           0001         16         17         18         19         20           0000         0         1         2         3         4 | 0000         0001         0010         0011         0100         0101           0111         112         113         114         115         116         117           0110         96         97         98         99         100         101           0101         80         81         82         83         84         85           0100         64         65         66         67         68         69           0011         48         49         50         51         52         53           0010         32         33         34         35         36         37           0001         16         17         18         19         20         21           0000         0         1         2         3         4         5 | 0000         0001         0010         0011         0100         0101         0110           0111         112         113         114         115         116         117         118           0110         96         97         98         99         100         101         102           0101         80         81         82         83         84         85         86           0100         64         65         66         67         68         69         70           0011         48         49         50         51         52         53         54           0010         32         33         34         35         36         37         38           0001         16         17         18         19         20         21         22           0000         0         1         2         3         4         5         6 | 0000         0001         0010         0011         0100         0101         0110         0111           0111         112         113         114         115         116         117         118         119           0110         96         97         98         99         100         101         102         103           0101         80         81         82         83         84         85         86         87           0100         64         65         66         67         68         69         70         71           0011         48         49         50         51         52         53         54         55           0010         32         33         34         35         36         37         38         39           0001         16         17         18         19         20         21         22         23           0000         0         1         2         3         4         5         6         7 |

# Board representation—0x88

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|---|-----------|------|------|------|------|------|------|------|------|-----|
|   |           | 0000 | 0001 | 0010 | 0011 | 0100 | 0101 | 0110 | 0111 | low |
| 8 | 0111      | 112  | 113  | 114  | 115  | 116  | 117  | 118  | 119  |     |
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| 6 | 0101      | 80   | 81   | 82   | 83   | 84   | 85   | 86   | 87   |     |
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| 2 | 0001      | 16   | 17   | 18   | 19   | 20   | 21   | 22   | 23   |     |
| 1 | 0000      | 0    | 1    | 2    | 3    | 4    | 5    | 6    | 7    |     |
|   | high bits |      |      |      |      |      |      |      |      |     |

To move one field to the left or right, just subtract or add one.

# Board representation—0x88

|   |           | a    | b    | С    | d    | е    | f    | g    | h    |     |
|---|-----------|------|------|------|------|------|------|------|------|-----|
|   |           | 0000 | 0001 | 0010 | 0011 | 0100 | 0101 | 0110 | 0111 | low |
| 8 | 0111      | 112  | 113  | 114  | 115  | 116  | 117  | 118  | 119  |     |
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| 3 | 0010      | 32   | 33   | 34   | 35   | 36   | 37   | 38   | 39   |     |
| 2 | 0001      | 16   | 17   | 18   | 19   | 20   | 21   | 22   | 23   |     |
| 1 | 0000      | 0    | 1    | 2    | 3    | 4    | 5    | 6    | 7    |     |
|   | high bits |      |      |      |      |      |      |      |      |     |

To move up a row, add 16, to move down a row, subtract 16.

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|---|-----------|------|------|------|------|------|------|------|------|-----|
|   |           | 0000 | 0001 | 0010 | 0011 | 0100 | 0101 | 0110 | 0111 | low |
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| 1 | 0000      | 0    | 1    | 2    | 3    | 4    | 5    | 6    | 7    |     |
|   | high bits |      |      |      |      |      |      |      |      |     |

Board: represented as an array with 128 entries, only 64 of which correspond to actual fields.

# **Board representation**-0x88

|   |           | a    | b    | С    | d    | е    | f    | g    | h    |     |
|---|-----------|------|------|------|------|------|------|------|------|-----|
|   |           | 0000 | 0001 | 0010 | 0011 | 0100 | 0101 | 0110 | 0111 | low |
| 8 | 0111      | 112  | 113  | 114  | 115  | 116  | 117  | 118  | 119  |     |
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| 1 | 0000      | 0    | 1    | 2    | 3    | 4    | 5    | 6    | 7    |     |
|   | high bits |      |      |      |      |      |      |      |      |     |

This is much faster than the first version. To check whether a number i is a valid position on the board, check whether it satisfies i&0x88 == 0 (&: bitwise).

Idea: for each colour and piece, use a 'bitboard'.

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The white pawns:



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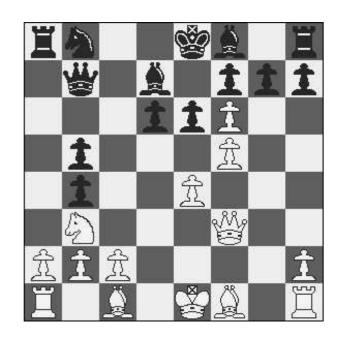
The white pawns:



| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
|---|---|---|---|---|---|---|---|
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 1 | 1 | 0 | 0 | 0 | 0 | 1 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

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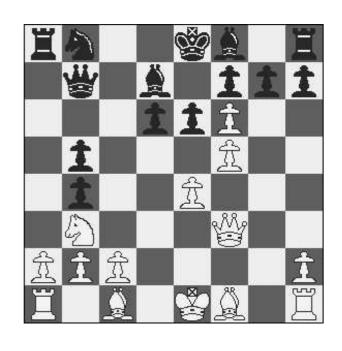


| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
|---|---|---|---|---|---|---|---|
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 1 | 1 | 0 | 0 | 0 | 0 | 1 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

Need: one 64-bit word for each piece. Operations: bit-wise—this is really fast!

Idea: for each colour and piece, use a 'bitboard'.





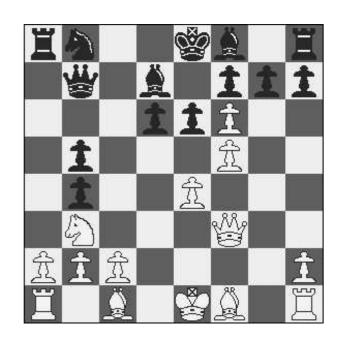
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
|---|---|---|---|---|---|---|---|
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 1 | 1 | 0 | 0 | 0 | 0 | 1 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

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Example: bitboard for all black pieces: bit-wise 'or' of all bitboards for black pieces.

Idea: for each colour and piece, use a 'bitboard'.





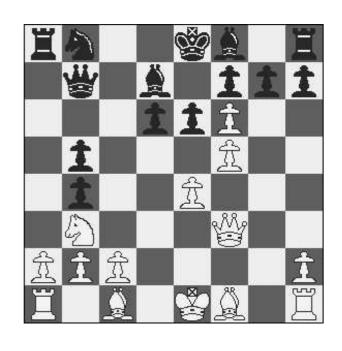
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
|---|---|---|---|---|---|---|---|
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 1 | 1 | 0 | 0 | 0 | 0 | 1 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

Need: one 64-bit word for each piece. Operations: bit-wise—this is really fast!

Move of a piece by a row: shift the bitboard by 8.

Idea: for each colour and piece, use a 'bitboard'.





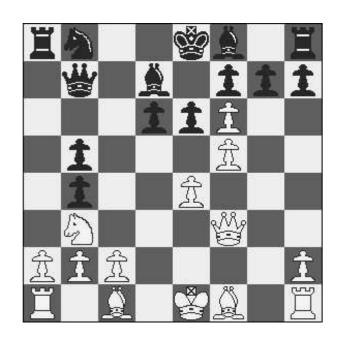
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
|---|---|---|---|---|---|---|---|
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 1 | 1 | 0 | 0 | 0 | 0 | 1 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

Need: one 64-bit word for each piece. Operations: bit-wise—this is really fast!

Empty fields: bitboard for all pieces negated.

Idea: for each colour and piece, use a 'bitboard'.

The white pawns:

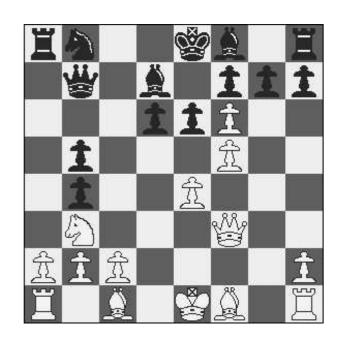


| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
|---|---|---|---|---|---|---|---|
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 1 | 1 | 0 | 0 | 0 | 0 | 1 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

All legal moves of pawns by one field can be stored in a bitboard (similarly for all legal moves of pawns by two fields). Constant bitboards can be prepared at compile time to be available in a library.

Idea: for each colour and piece, use a 'bitboard'.

The white pawns:

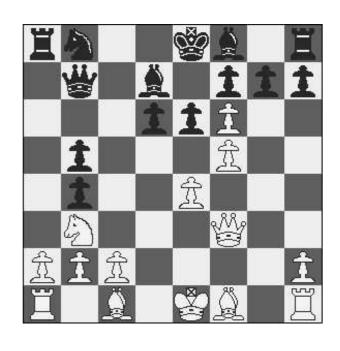


| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
|---|---|---|---|---|---|---|---|
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 1 | 1 | 0 | 0 | 0 | 0 | 1 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

Pawn captures: shifting the bitboard by 7 or 9 and bit-wise 'and' it with the bitboard for pieces of the opposite colour.

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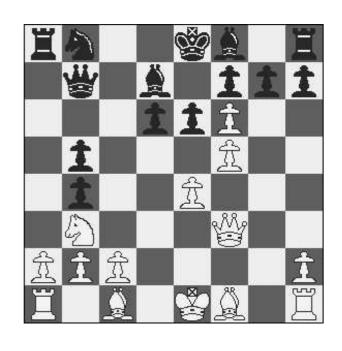


| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
|---|---|---|---|---|---|---|---|
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 1 | 1 | 0 | 0 | 0 | 0 | 1 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

Only disadvantage: the code becomes more complicated; turning a bitboard of possible moves into a list of possible moves, for example.

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| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
|---|---|---|---|---|---|---|---|
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 1 | 1 | 0 | 0 | 0 | 0 | 1 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

Only disadvantage: the code becomes more complicated; turning a bitboard of possible moves into a list of possible moves, for example.

Advantages: fast; bitboards required more than once only have to be computed once; several moves can be generated at the same time.

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A hash function frequently used consists of assigning to each pair, consisting of a piece and a field on the board, a large random number. The idea is that this number encodes the fact that the corresponding piece occupies the corresponding field. Then one sums up the appropriate numbers for the given position to obtain the hash key. A checksum process can be applied to make sure later that 'the right' position is looked up.

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This is best done by keeping a stack of moves with sufficient information to undo them. This is typically much cheaper than keeping a list of positions through which one has gone.



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If the values provided do not judge the given position accurately then the program can't possibly play well—until very close to the end, they are all the program has to judge which move it should make.

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There are no hard and fast rules for what makes a good evaluation function; they are mostly based on heuristics.

When writing a game-playing program, speed is always an issue. Hence it pays to calculate the desired evaluation function in such a way to make the process as fast as possible.

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Let p be the current position, and e the evaluation function. Then if

$$e(p) = e_{s_1}(s_1)$$
's place in  $p + \cdots + e_{s_n}(s_n)$ 's place in  $p + \cdots + e_{s_n}(s_n)$ 's place in  $p + \cdots + e_{s_n}(s_n)$ 

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where  $s_1, \ldots, s_n$  are the pieces involved, the value of a new position resulting from one piece s being moved is

$$score(move) = e_s(s's new field) - e_s(s's old field).$$

Problems: For many games this kind of evaluation function is not good enough—it does not take the relative position of pieces into account.

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It is important that an evaluation function judge any position from both players' point of view. Having many pieces on the board does not give White any advantage if Black is about to checkmate him!

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Not equally useful for all games: Othello: not number of pieces is important, but their locations (corners). Player with fewer pieces might have better position. There are other games where the number of pieces may be irrelevant.

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Chess: count number of fields threatened by one player; Othello: count number of pieces which cannot be taken by the opponent. Calculate size, possible with weights for very important fields.

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Mobility. Ability to move. Having many different available moves: advantageous, *e.g.* in Othello. Chess: not clear this is useful.

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Other games: try 'parity argument': often find positions where player who moves next wins/loses, can be simple to evaluate (see Nim, Connect-4).

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Threats. Can one of the players capture (or threaten to capture) a piece? Connect-4, Go-Moku: can a player win in the next move? Othello: is a player threatening to take a corner?

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Shape. How pieces on the board relate to each other. Chess: line of pawns much stronger than other grouping. Go: shape is 'territory to be'—stones outline territory which the player can defend when threatened.

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Judging shape: often very difficult. Change of shape value: incremental over time, long-term target. Evaluation function partially based on shape: can't just simply add piece-based functions.

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Mobility. Ability to move.

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Known Patterns. Go: libraries of sequences of moves in small areas (joseki)—preserves balance between players.

#### Relevant constituent parts

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Mobility. Ability to move.

Tempo. Initiative.

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Chess: bishop capturing a pawn on border is often trapped; Othello: sacrifice one corners in exchange for another. Deciding when a pattern applies is hard!

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Know, *e.g.*: one rook less than two pawns and bishop, or two pawns and knight, but not less than one pawn and bishop/knight.

So: weight of a rook should be below weight of pawns and bishop, but above one pawn and bishop. Get fewer possibilities to try.

#### Deducing constraints.

Hand tweaking. Happens typically in practice. Programmers watch implementation play, judge which parameters to change and how. Perform the change and watch again. Reasonably fast but requires game-specific knowledge.

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Can be modified by randomly sticking with some changes which do not improve performance. 'Randomness' controlled by some probabilities (start out fairly high, become smaller as a good value is approached). Adjusted method is slower than original, but can get good values.

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Examples for learning: genetic algorithms, neural networks. Both: rather slow; main advantage: do not require game-specific knowledge. Reason for slowness: number of test games required is typically very high (commercial game programmers tried about 3000 matches to allow the program to learn—the result was worse than hand tweaking). Further problem: If opponent is too good program loses all the time and never starts learning.

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Problem with playing program against versions of itself: same lines are explored over and over. To avoid this: start the program(s) from positions a few moves into a game.



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There are some ways of fiddling with this to adjust it to the game in question. The thought is always to make it faster so that it can search deeper.

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Obvious advantage: When time runs out we give the best move found so far, and that will at least be sensible. This is known as iterative deepening.

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But if we use a hash table to keep track of results so far we can estimate a value.

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- successively increasing lower bounds for a max node  $(\alpha)$ ;
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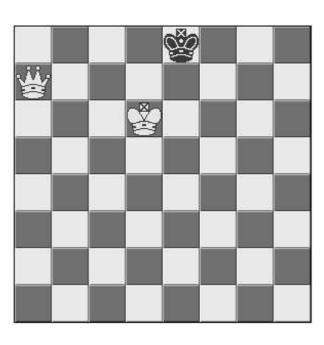
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Can search the first move(s) with big window for potential value (see aspiration search), and later moves with smaller ones. This is known as principal variation search.

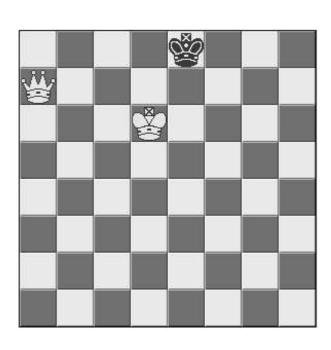




But:

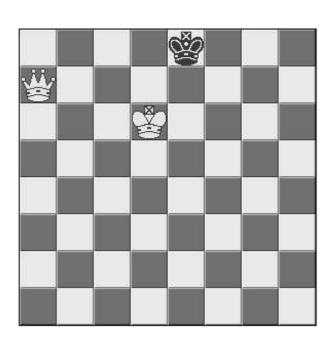


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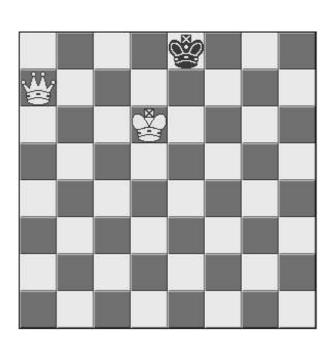
If White moves the king to e6 (one field to the right) then he is still in a winning position, with Black's only valid moves being to d8 and f8.

But:



Let's assume Black moves to d8. Then moving the king back to d6 again gives White a winning position.

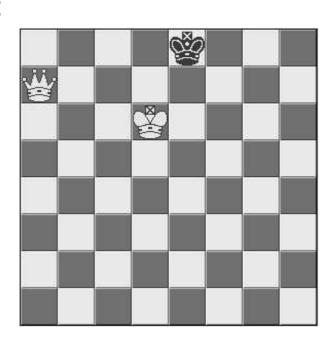
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Can avoid this by assigning slightly lower values to winning positions, for example

1000 – number of moves req'd to get win.

Then alpha-beta search will work properly.

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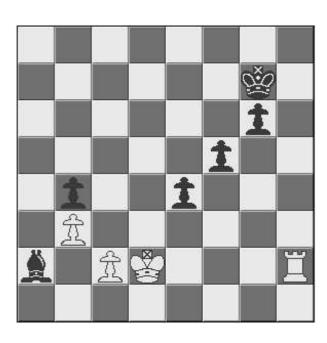
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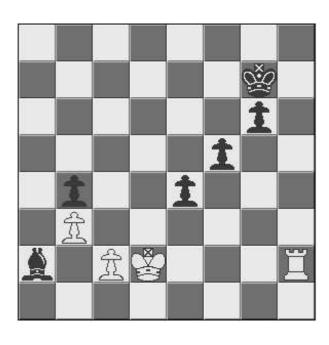
This can lead to the program trying to fend off bad events (capture of its piece, for example) by keeping them below the horizon.

In order to avoid, say, the capture of one of its pieces the program may try pointless moves which merely postpone the inevitable—typically these moves do not progress the program's play.

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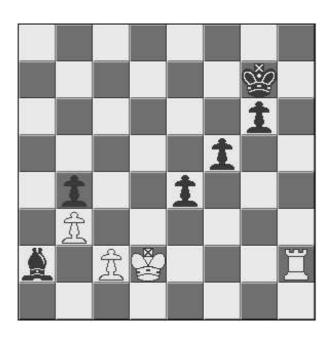


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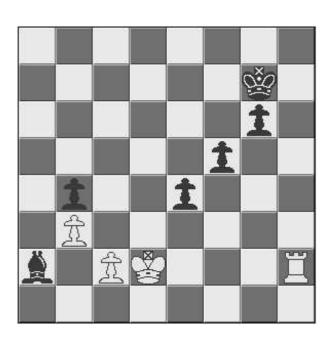


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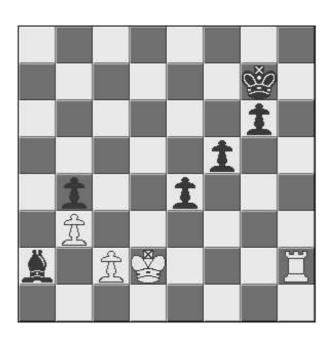
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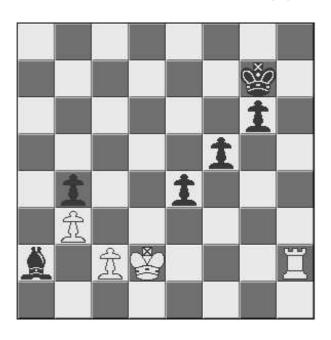
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Solutions: Add knowledge so that program can detect when piece is trapped. Increase overall depth of search in such situations so that horizon is widened. Whenever piece is threatened, search to deeper level selectively.

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Many programs search deeper on what they think is the best move (see principal variation search).

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Mistakes or weaknesses in a program can be explored over and over (until the creator finds a chance to fix this, since these programs don't learn). Many tournaments between various programs seemed to be more about who could discover whose built-in faults, rather than whose program genuinely played best!



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Shannon thought this would be a useful application for computers, and would give insights into how one makes intelligent decisions.

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1974: First world computer Chess championships. Repeated every three years.

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Search to variable depth, depending on whether the current position is judged to be 'tricky' or relatively straight-forward.

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The early, strong claims regarding the possibilities of AI turned to out to be vastly exaggerated. Today, Artificial Intelligence often is about search techniques and the machine learning is very different from human learning!

## **Further improvements**

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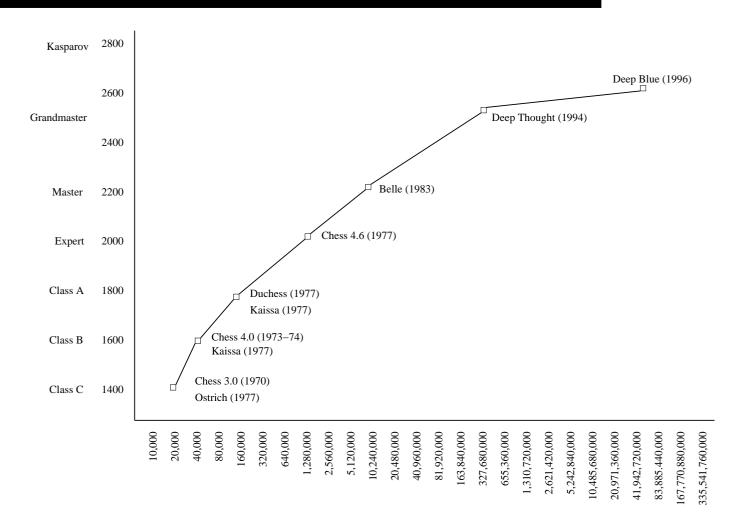
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Since late eighties: Main development has gone into specialized hardware.

# Speed increases strength



Number of positions examined in three minutes, official ranking. (Note logarithmic scale along horizontal axis!) Where is perfect play?

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|   | 3  | 4    | 5    | 6    | 7    | 8   | rating |
|---|----|------|------|------|------|-----|--------|
| 3 |    | 4    |      |      |      |     | 1091   |
| 4 | 16 |      | 5.5  |      |      |     | 1332   |
| 5 |    | 14.5 |      | 4.5  |      |     | 1500   |
| 6 |    |      | 15.5 |      | 2.5  |     | 1714   |
| 7 |    |      |      | 17.5 |      | 3.5 | 2052   |
| 8 |    |      |      |      | 16.5 |     | 2320   |

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|   | 4    | 5    | 6    | 7  | 8    | 9   | rating |
|---|------|------|------|----|------|-----|--------|
| 4 |      | 5    | .5   | 0  | 0    | 0   | 1235   |
| 5 | 15   |      | 3.5  | 3  | .5   | 0   | 1570   |
| 6 | 19.5 | 16.5 |      | 4  | 1.5  | 1.5 | 1826   |
| 7 | 20   | 17   | 16   |    | 5    | 4   | 2031   |
| 8 | 20   | 19.5 | 18.5 | 15 |      | 5.5 | 2208   |
| 9 | 20   | 20   | 18.5 | 16 | 14.5 |     | 2328   |

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Three or four levels more of search means outclassing one's opponent!

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|       | percent. of moves |        |
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| level | from predecessor  | rating |
| 4     | 33.1              | 1300   |
| 5     | 33.1              | 1570   |
| 6     | 27.7              | 1796   |
| 7     | 29.5              | 2037   |
| 8     | 26.0              | 2249   |
| 9     | 22.6              | 2433   |
| 10    | 17.7              | 2577   |
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| level | from predecessor  | rating | are picked when search-                 |
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### **Hardware for Chess**

The following table gives an overview over Chess-playing programs and the hardware they were running on.

### **Hardware for Chess**

| Name             | Year | Description                           |
|------------------|------|---------------------------------------|
| Ostrich          | 1981 | 5-processor Data General system       |
| Ostrich          | 1982 | 8-processor Data General system       |
| Cray Blitz       | 1983 | 2-processor Cray XMP                  |
| Cray Blitz       | 1984 | 4-processor Cray XMP                  |
| Sun Phoenix      | 1986 | Network of 20 VAXs and Suns           |
| Chess Challenger | 1986 | 20 8086 microprocessors               |
| Waycool          | 1986 | 64-processor N/Cube system            |
| Waycool          | 1988 | 256-processor N/Cube system           |
| Deep Thought     | 1989 | 3 2-processor VLSI chess circuits     |
| Star Tech        | 1993 | 512-processor Connection Machine      |
| Star Socrates    | 1995 | 1,824-processor Intel Paragon         |
| Zugzwang         | 1995 | 96-processor GC-Powerplus distributed |
|                  |      | system (based on the PowerPC)         |
| Deep Blue        | 1996 | 32-processor IBM RS/6000 SP with      |
|                  |      | 6 VLSI chess circuits per processor   |

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1997: Rematch, ending 2.5 to 3.5, Kasparov makes mistake in final and deciding match.

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- Board representations should make the generation of moves, doing and undoing them fast.
- Evaluation functions assign a value to a board position statically, from just what is on the board. They require knowledge about the game in question.

- Three tasks have to be solved: Designing an internal board representation and generating valid moves, designing an evaluation function and implementing (some variant of) alpha-beta search. All considerations are overshadowed by the need for speed.
- Board representations should make the generation of moves, doing and undoing them fast.
- Evaluation functions assign a value to a board position statically, from just what is on the board.
- Alpha-beta search assigns a value to a position by searching the game tree below it and eventually applying the evaluation function. Searching to greater depth will result in a better program, so any gain in speed goes into searching to a greater depth. There are many tricks to try to only search the relevant parts of the game tree; in particular ordering moves to search the most promising ones first.

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- Most effort so far has gone into creating Chess-playing programs. They have profited from faster hardware, and many improvements have been made which are very Chess-specific: better heuristics, opening and endgame libraries, and the like.