DAML+OIL and Description Logic Reasoning

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Talk Outline

The Semantic Web and DAML+OIL

Description Logics and Reasoning

Reasoning techniques

Implementing DL systems

Research Challenges

Summary

The Semantic Web and DAML+OIL

Semantic Web Ontology Languages

US **DAML** programme (in cooperation with W3C and a cast of thousands) aim to develop so-called **Semantic Web**

- Most existing Web resources only human understandable
 - Markup (HTML) provides rendering information
 - Textual/graphical information for human consumption
- Semantic Web aims at machine understandability
 - Semantic markup will be added to web resources
 - Markup will use Ontologies for shared understanding
- Requirement for a suitable ontology language
 - Compatible with existing Web standards (XML, RDF, RDFS)
 - Captures common KR idioms
 - Formally specified and of adequate expressive power
 - Can provide reasoning support
- DAML-ONT language developed to meet these requirements

OIL and DAML+OIL

Meanwhile, somewhere in darkest Europe...

- OIL language already developed to meet similar requirements
 - Extends existing Web standards (XML, RDF, RDFS)
 - Intuitive (frame) syntax plus high expressive power
 - Well defined semantics via mapping to SHIQ DL
 - Can use DL systems to reason with OIL ontologies
- Two efforts merged to produce single language, DAML+OIL
- Detailed specification agreed by Joint EU/US Committee on Agent Markup Languages
- Proposed W3C Ontology Language WG will take DAML+OIL as starting point (?)

DAML+OIL Language Overview

DAML+OIL is an ontology language

- Describes structure of the domain (i.e., a Tbox)
 - RDF used to describe specific instances (i.e., an Abox)
- Structure described in terms of classes (concepts) and properties (roles)
- Ontology consists of set of axioms
 - E.g., asserting class subsumption/equivalence
- Classes can be names or expressions
 - Various constructors provided for building class expressions
- Expressive power determined by
 - Kinds of axiom supported
 - Kinds of class (and property) constructor supported

DAML+OIL Overview: Class Constructors

Constructor	DL Syntax	Example
intersectionOf	$C_1 \sqcap \ldots \sqcap C_n$	Human □ Male
unionOf	$C_1 \sqcup \ldots \sqcup C_n$	Doctor ⊔ Lawyer
complementOf	$\neg C$	¬Male
oneOf	$\{x_1 \dots x_n\}$	{john, mary}
toClass	$\forall P.C$	∀hasChild.Doctor
hasClass	$\exists P.C$	∃hasChild.Lawyer
hasValue	$\exists P.\{x\}$	∃citizenOf.{USA}
minCardinalityQ	$\geqslant nP.C$	≥2hasChild.Lawyer
maxCardinalityQ	$\leq nP.C$	≼1hasChild.Male
cardinalityQ	=n P.C	=1 hasParent.Female

- XMLS datatypes as well as classes
- Arbitrarily complex nesting of constructors
 - E.g., ∀hasChild.(Doctor

 ∃hasChild.Doctor)

DAML+OIL Overview: Axioms

Axiom	DL Syntax	Example
subClassOf	$C_1 \sqsubseteq C_2$	Human ⊑ Animal ⊓ Biped
sameClassAs	$C_1 \doteq C_2$	Man ≐ Human ⊓ Male
subPropertyOf	$P_1 \sqsubseteq P_2$	hasDaughter ⊑ hasChild
samePropertyAs	$P_1 \doteq P_2$	cost ≐ price
sameIndividualAs	$\begin{cases} x_1 \rbrace \doteq \{x_2 \rbrace \end{cases}$	$\{President_Bush\} \doteq \{G_W_Bush\}$
disjointWith	$C_1 \sqsubseteq \neg C_2$	Male ⊑ ¬Female
differentIndividualFrom	$ \{x_1\} \sqsubseteq \neg \{x_2\} $	$\{john\} \sqsubseteq \neg \{peter\}$
inverseOf	$P_1 \doteq P_2^-$	$hasChild \doteq hasParent^-$
transitiveProperty	$P^+ \sqsubseteq P$	ancestor ⁺ ⊑ ancestor
uniqueProperty	$\top \sqsubseteq \leqslant 1P$	$ op \sqsubseteq \leqslant 1$ hasMother
unambiguousProperty	$\top \sqsubseteq \leqslant 1P^-$	$ op \sqsubseteq \leqslant 1$ isMotherOf $^-$

Axioms (mostly) reducible to subClass/PropertyOf

DAML+OIL

- Is a Description Logic (but don't tell anyone)
- ightharpoonup More precisely, DAML+OIL is SHIQ
 - Plus nominals
 - Plus datatypes (simple concrete domains)
 - With RDFS based syntax
- - SHIQ is based on 15+ years of DL research
- Can use DL reasoning with DAML+OIL
 - Existing SHIQ implementations support (most of) DAML+OIL

Why Reasoning Services?

Reasoning is important for:

- Ontology design
 - Check class consistency and (unexpected) implied relationships
 - Particularly important with large ontologies/multiple authors
- Ontology integration
 - Assert inter-ontology relationships
 - Reasoner computes integrated class hierarchy/consistency
- Ontology deployment
 - Determine if set of facts are consistent w.r.t. ontology
 - Determine if individuals are instances of ontology classes

"The Semantic Web needs a logic on top" (Henry Thompson)

Why Decidable Reasoning?

Set of operators/axioms restricted so that reasoning is decidable

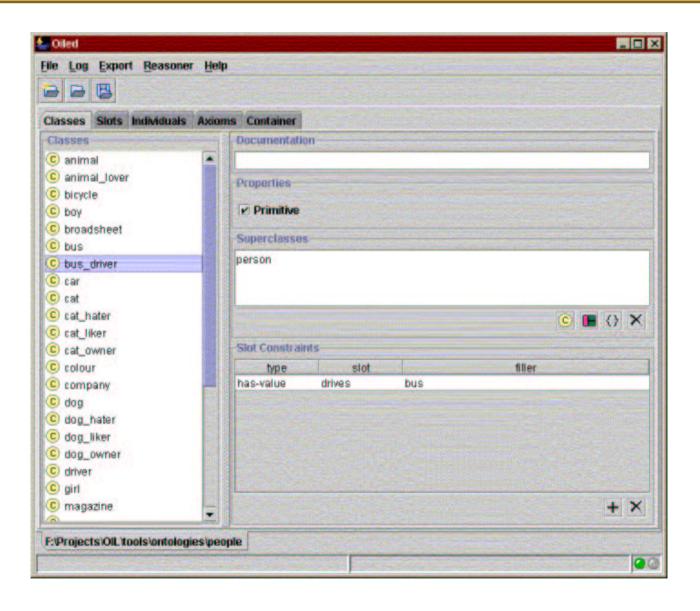
- Consistent with Semantic Web's layered architecture
 - XML provides syntax transport layer
 - RDF provides basic relational language
 - RDFS provides basic ontological primitives
 - DAML+OIL provides (decidable) logical layer
 - Further layers (e.g., rules) will extend DAML+OIL
 - Extensions will almost certainly be undecidable
- Facilitates provision of reasoning services
 - Known algorithms
 - Implemented systems
 - Evidence of empirical tractability

Reasoning Support for Ontology Design: OilEd

OilEd is a DAML+OIL ontology editor with DL reasoning support

- Frame based interface (inspired by Protégé)
 - Classes defined by superclass(es) plus slot constraints
- Extended to clarify semantics and capture whole language
 - Primitive (□) and defined (=) classes
 - Explicit ∃ (hasClass), ∀ (toClass) and cardinality restrictions
 - Boolean connectives (□, □, ¬) and nesting
 - Transitive, symmetrical and functional properties
 - Disjointness, inclusion (⊆) and equality (≐) axioms
 - Fake individuals
- Reasoning support provided by FaCT system
 - Ontology translated into SHIQ DL
 - Communicates with FaCT via CORBA interface
 - Indicates inconsistencies and implicit subsumptions

OilEd



Description Logics and Reasoning

What are Description Logics?

- Based on concepts (classes) and roles
 - Concepts (classes) are interpreted as sets of objects
 - Roles are interpreted as binary relations on objects
- Descendants of semantic networks and KL-ONE
- Decidable fragments of FOL
 - Many DLs are fragments of L2, C2 or the Guarded Fragment
- Closely related to propositional modal logics
- Also known as terminological logics, concept languages, etc.
- Key features of DLs are
 - Well defined semantics (they are logics)
 - Provision of inference services

Short History of Description Logics

Phase 1:

- Incomplete systems (Back, Classic, Loom, ...)
- Based on structural algorithms

Phase 2:

- Development of tableau algorithms and complexity results
- Tableau-based systems (Kris, Crack)
- Investigation of optimisation techniques

Phase 3:

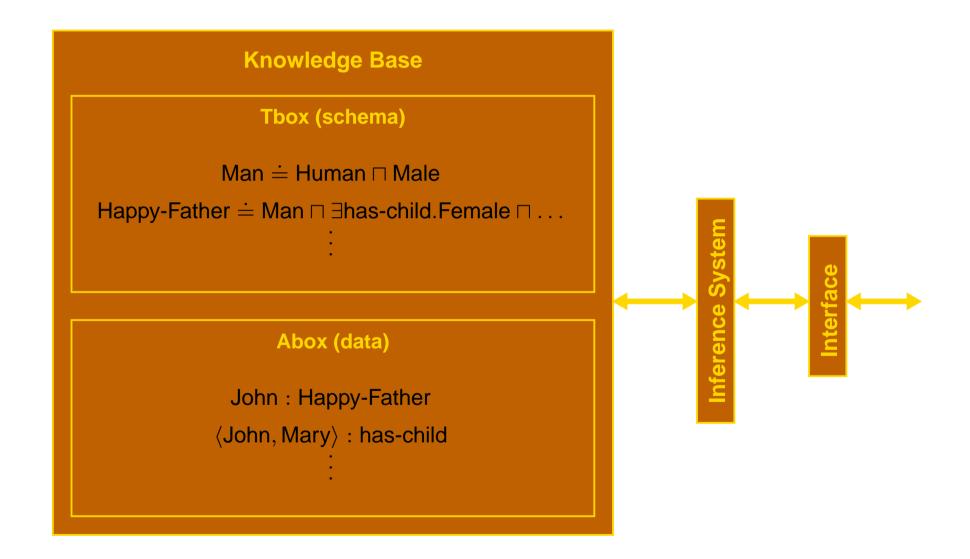
- Tableau algorithms for very expressive DLs
- Highly optimised tableau systems (FaCT, DLP, Racer)
- Relationship to modal logic and decidable fragments of FOL

Latest Developments

Phase 4:

- Mature implementations
- Mainstream applications and Tools
 - Databases
 - Consistency of conceptual schemata
 - Schema integration
 - Query subsumption (w.r.t. a conceptual schema)
 - Ontologies and Semantic Web
 - Design and Maintenance
 - Integration
 - Deployment
- Commercial implementations
 - Cerebra system from Network Inference Ltd

DL System Architecture



DL Constructors

Particular DLs characterised by **set of constructors** provided for building complex concepts and roles from simpler ones

- Usually include at least:
 - Conjunction (□), disjunction (□), negation (¬)
 - Restricted (guarded) forms of quantification (∃, ∀)
- This basic DL is known as ALC

DL Syntax and Semantics

Semantics given by interpretation $\mathcal{I} = (\Delta^{\mathcal{I}}, \cdot^{\mathcal{I}})$

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Constructor	Syntax	Example	Semantics	
atomic concept	A	Human	$A^{\mathcal{I}} \subseteq \Delta^{\mathcal{I}}$	
atomic role	R	has-child	$R^{\mathcal{I}} \subseteq \Delta^{\mathcal{I}} \times \Delta^{\mathcal{I}}$	
and for C , D concepts and R a role name				
conjunction	$C\sqcap D$	Human ⊓ Male	$C^{\mathcal{I}} \cap D^{\mathcal{I}}$	
disjunction	$C \sqcup D$	Doctor ⊔ Lawyer	$C^{\mathcal{I}} \cup D^{\mathcal{I}}$	
negation	$\neg C$	⊣Male	$\Delta^{\mathcal{I}} \setminus C$	
exists restr.	$\exists R.C$	∃has-child.Male	$\{x \mid \exists y. \langle x, y \rangle \in R^{\mathcal{I}} \land y \in C^{\mathcal{I}}\}$	
value restr.	$\forall R.C$	∀has-child.Doctor		

Other DL Constructors

Many different DLs/DL constructors have been investigated, e.g.

Constructor	Syntax	Example	Semantics
number restr.	$\geqslant nR$	≽3 has-child	$\{x \mid \{y.\langle x, y\rangle \in R^{\mathcal{I}}\} \geqslant n\}$
	$\leq nR$	\leqslant 1 has-mother	$\{x \mid \{y.\langle x,y\rangle \in R^{\mathcal{I}}\} \leqslant n\}$
inverse role	R^{-}	has-child ⁻	$\{\langle x, y \rangle \mid \langle y, x \rangle \in R^{\mathcal{I}}\}$
trans. role	R^*	has-child*	$(R^{\mathcal{I}})^*$
concrete domain	$f_1,\ldots,f_n.P$	earns spends <	$\{x \mid P(f_1^{\mathcal{I}}, \dots, f_n^{\mathcal{I}})\}$

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DL Knowledge Base (Tbox)

Terminological part (Tbox) is set of axioms describing structure of domain

Definition axioms introduce macros/names for concepts

$$A \doteq C$$
, $A \sqsubseteq C$

Human

☐ Animal
☐ Biped

Inclusion (GCI) axioms assert subsumption relations

 $C \sqsubseteq D$ (note $C \doteq D$ equivalent to $C \sqsubseteq D$ and $D \sqsubseteq C$)

∃has-degree.Masters ⊑ ∃has-degree.Bachelors

DL Knowledge Base (Abox)

Assertional part (Abox) is set of axioms describing concrete situation

Concept assertions

a:C

John : Man

∃has-child.Female

Role assertions

 $\langle a,b\rangle:R$

⟨John, Mary⟩ : has-child

Basic Inference Problems

Subsumption (structure knowledge, compute taxonomy)

 $C \sqsubseteq D$? Is $C^{\mathcal{I}} \subseteq D^{\mathcal{I}}$ in all interpretations?

Subsumption w.r.t. Tbox T

 $C \sqsubseteq_{\mathcal{T}} D$? Is $C^{\mathcal{I}} \subseteq D^{\mathcal{I}}$ in all models of \mathcal{T} ?

Consistency

Is C consistent w.r.t. \mathcal{T} ? Is there a model \mathcal{I} of \mathcal{T} s.t. $C^{\mathcal{I}} \neq \emptyset$?

KB Consistency

Is $\langle \mathcal{T}, \mathcal{A} \rangle$ consistent? Is there a model \mathcal{I} of $\langle \mathcal{T}, \mathcal{A} \rangle$?

Reasoning Techniques

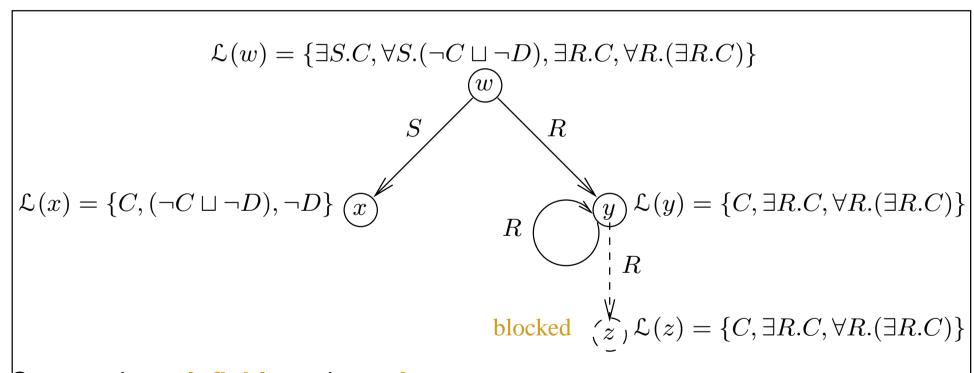
Subsumption and Satisfiability

Subsumption transformed into satisfiability Tableaux algorithm used to test satisfiability

- Try to build model (witness) of concept C
- Model represented by tree T
 - Nodes in T correspond to individuals in model
 - Nodes labeled with sets of subconcepts of C
 - Edges labeled with role names in C
- rightharpoonup Start from root node labeled $\{C\}$
- Apply expansion rules to node labels until
 - Rules correspond with language constructs
 - Expansion completed (tree represents valid model)
 - Contradictions prove there is no model
- \longrightarrow Non-deterministic expansion \longrightarrow search (e.g., $C \sqcup D$)
- Blocking ensures termination (with expressive DLs)

Tableaux Expansion

Test satisfiability of $\exists S.C \sqcap \forall S.(\neg C \sqcup \neg D) \sqcap \exists R.C \sqcap \forall R.(\exists R.C)\}$ where R is a transitive role



Concept is satisfiable: w is a witness

More Advanced Techniques

Satisfiability w.r.t. a Terminology

For each GCI $C \sqsubseteq D \in \mathcal{T}$, add $\neg C \sqcup D$ to every node label

More expressive DLs

- Basic technique can be extended to deal with
 - Role inclusion axioms (role hierarchy)
 - Number restrictions
 - Inverse roles
 - Concrete domains
 - Aboxes
 - etc.
- Extend expansion rules and use more sophisticated blocking strategy
- Forest instead of Tree (for Aboxes)

Implementing DL Systems

Naive Implementations

Problems include:

- Space usage
 - Storage required for tableaux datastructures
 - Rarely a serious problem in practice
 - But problems can arise with inverse roles and cyclical KBs
- Time usage
 - Search required due to non-deterministic expansion
 - Serious problem in practice
 - Mitigated by:
 - Careful choice of algorithm
 - → Highly optimised implementation

Careful Choice of Algorithm

- Transitive roles instead of transitive closure
 - Deterministic expansion of $\exists R.C$, even when $R \in \mathbf{R}_+$
 - (Relatively) simple blocking conditions
 - Cycles always represent (part of) valid cyclical models
- Direct algorithm/implementation instead of encodings
 - GCI axioms can be used to "encode" additional operators/axioms
 - Powerful technique, particularly when used with FL closure
 - Can encode cardinality constraints, inverse roles, range/domain,

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→ E.g., (domain R.C) $\equiv \exists R. \top \sqsubseteq C$

- (FL) encodings introduce (large numbers of) axioms
- BUT even simple domain encoding is disastrous with large numbers of roles

Highly Optimised Implementation

Modern systems include MANY optimisations, e.g.:

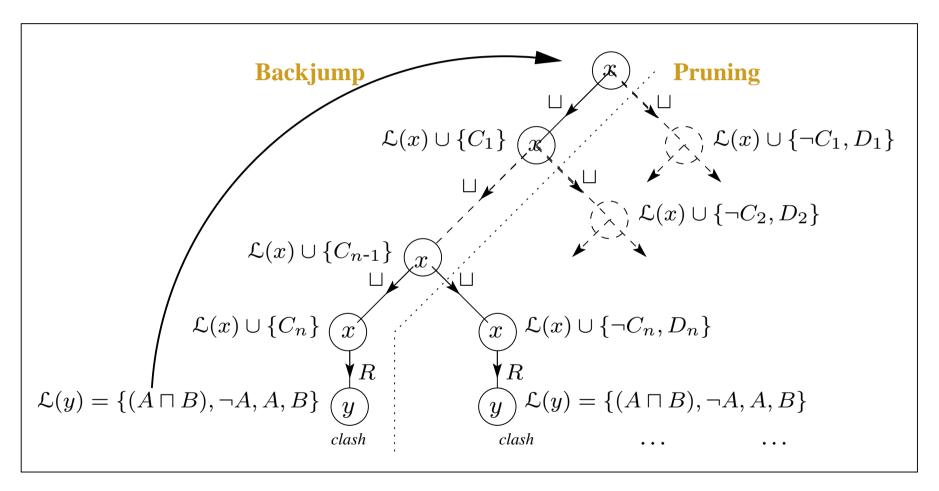
- Optimised classification
 - Use enhanced traversal (exploit information from previous tests)
 - Use structural information to select classification order
- Optimised subsumption testing
 - Normalisation and simplification of concepts
 - Absorption (simplification) of general axioms
 - Davis-Putnam style semantic branching search
 - Dependency directed backtracking
 - Caching
 - Heuristic ordering of propositional and modal expansion

Dependency Directed Backtracking

- Allows rapid recovery from bad branching choices
- Most commonly used technique is backjumping
 - Tag concepts introduced at branch points (e.g., when expanding disjunctions)
 - Expansion rules combine and propagate tags
 - On discovering a clash, identify most recently introduced concepts involved
 - Jump back to relevant branch points without exploring alternative branches
 - Effect is to prune away part of the search space
- Highly effective essential for usable system
 - E.g., GALEN KB, 30s (with) → months++ (without)

Backjumping

E.g., if $\exists R. \neg A \sqcap \forall R. (A \sqcap B) \sqcap (C_1 \sqcup D_1) \sqcap \ldots \sqcap (C_n \sqcup D_n) \subseteq \mathcal{L}(x)$



Research Challenges

Research Challenges

Increased expressive power

- Datatypes
- Nominals
- Extensions to DAML+OIL

Performance

- Inverse roles and qualified number restrictions
- Very large KBs
- Reasoning with individuals

Tools and Infrastructure

Support for large scale ontological engineering and deployment

New reasoning tasks

- Querying
- Lcs/matching
- Sanctioning

• . . .

Increased Expressive Power: Datatypes

DAML+OIL extends SHIQ with datatypes and nominals

Datatypes

- DAML+OIL has simple form of datatypes
 - Unary predicates plus disjoint abstract/datatype domains
- Theoretically not particularly challenging
 - Existing work on concrete domains [Baader & Hanschke, Lutz]
 - Algorithm already known for SHOQ(D) [Horrocks & Sattler]
- May be practically challenging
 - All XMLS datatypes supported
- Already seeing some (limited) implementations
 - E.g., Cerebra system (Network Inference)

Increased Expressive Power: Nominals

Nominals

- DAML+OIL has oneOf constructor
 - Extensionally defined concepts, e.g., $\{Mary\}^{\mathcal{I}} = \{Mary\}$
 - Equivalent to nominals in modal logic
- Theoretically very challenging
 - Resulting logic has known high complexity (NExpTime)
 - No known "practical" algorithm
 - Not obvious how to extend tableax techniques in this direction
 - Loss of tree model property
 - \rightarrow Spy-points: $\top \sqsubseteq \exists R. \{Spy\}$
 - → Finite domains: $\{Spy\} \sqsubseteq \leqslant nR^-$
- Relatively straightforward (in theory) without inverse roles
 - Algorithm for $\mathcal{SHOQ}(\mathbf{D})$ deals with nominals
 - Practical implementation still to be demonstrated

Increased Expressive Power: Extensions

- DAML+OIL not expressive enough for all applications
- Extensions wish list includes:
 - Feature chain (path) agreement, e.g., output of component of composite process equals input of subsequent process
 - Complex roles/role inclusions, e.g., a city located in part of a country is located in that country
 - Rules—proposal(s) already exist for "datalog/LP style rules"
 - Temporal and spatial reasoning
 - •
- May be impossible/undesirable to resist such extensions
- Extended language sure to be undecidable
- How can extensions best be integrated with DAML+OIL?
- How can reasoners be developed/adapted for extended languages
 - Some existing work on language fusions and hybrid reasoners

Performance Problems

- Evidence of empirical tractability mostly w.r.t. SHF— problems can arise when systems extended to SHIQ
- Important optimisations no longer (fully) work
 - E.g., problems with caching as cached models can affect parent
- Qualified number restrictions can also cause problems
 - Even relatively small numbers can mean significant non-determinism
- Reasoning with very large KBs/ontologies
 - Web ontologies can be expected to grow very large
- Reasoning with individuals (Abox)
 - Deployment of web ontologies will mean reasoning with (possibly very large numbers of) individuals
 - Unlikely that standard Abox techniques will be able to cope

Performance Solutions (Maybe)

Excessive memory usage

- Problem exacerbated by over-cautious double blocking condition (e.g., root node can never block)
- Promising results from more precise blocking condition [Sattler & Horrocks]

Qualified number restrictions

- Problem exacerbated by naive expansion rules
- Promising results from optimised expansion using Algebraic Methods [Haarslev & Möller]

Caching and merging

- Can still work in some situations (work in progress)
- Reasoning with very large KBs
 - DL systems shown to work with ≈100k concept KB [Haarslev & Möller]
 - But KB only exploited small part of DL language

Tools and Infrastructure

Tools and infrastructure required in order support use of DAML+OIL

- Ontology design and maintenance
 - Several editors available, e.g, OilEd (Manchester), OntoEdit (Karlsruhe), Protégé (Stanford)
 - Need integrated environments including modularity, versioning, visualisation, explanation, high-level languages, . . .
- Ontology Integration
 - Some tools available, e.g., Chimera (Stanford)
 - Need integrated environments . . .
 - Can learn from DB integration work [Lenzerini, Calvanese et al]
- Reasoning engines
 - Several DL systems available
 - Need for improved usability/connectivity

New Reasoning Tasks

Querying

- Retrieval (instances of a concept) and realisation (most specific class of instance) wont be sufficient
- Minimum requirement will be conjunctive query style language [Tessaris & Horrocks]
- May also need to answer "what can I say about x?" style of query [Bechhofer & Horrocks]
- Explanation (e.g., to support ontology design) [McGuinness, Borgida et al]
- Least common subsumer and/or matching (e.g., to support ontology integration and "bottom up" design) [Baader, Küsters & Molitor]

Summary

- Ontologies will play key role in Semantic Web
- DAML+OIL is web ontology language based on Description Logic
- Ontology design, integration and deployment supported by reasoning
- DLs are logic based KR formalisms with emphasis on reasoning
- DL systems provide efficient reasoning services
 - Careful choice of logic/algorithm
 - Highly optimised implementation
- Still many challenges for DL and Semantic Web research
 - Expressive power
 - Performance
 - Tools and infrastructure
 - New reasoning tasks

Resources

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Slides from this talk
 www.cs.man.ac.uk/~horrocks/Slides/hp-labs.pdf
FaCT system
 www.cs.man.ac.uk/fact
OIL
 www.ontoknowledge.org/oil/
DAML+OIL
 www.daml.org/language/
OilEd
 img.cs.man.ac.uk/oil
I.COM
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