# CS3191 Section 1 <br> Introduction to Games 

Andrea Schalk

Department of Computer Science, University of Manchester

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- Games and evolution. Games used to model the evolution of traits.


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- a game tree (required for the formal definition of game);
- a strategy;
- a pay-off function;
- the normal form of a game.


## What is a game?

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A game gives the rules for the interactions of the participating entities, the players.

## Example



## A game tree



The nodes in such a tree are the positions of the game.

## A game tree



At each node it has to be clear whose turn it is.

## A game tree



The bottom nodes in the tree, the leaves, are the final positions of the game. This is where we note the result.

## A game tree



Note that different positions in the tree can correspond to the same constellation on the board.

## A game tree



That is because in a tree, given a position we can reconstruct the path from the root leading to it (and thus the play of the game). Another path is shown. Taking trees instead of graphs makes our lives easier.

## Questions about game trees

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(b) Could you draw a game tree for Paper-Stone-Scissors? If not, why not?
(c) Consider the following simple game between two players: Player 1 has a coin which he hides under his hand, having first decided whether it should show head or tail. Player 2 guesses which of these has been chosen. If she guesses correctly, Player 1 pays her 1 quid, otherwise she has to pay the same amount to him. Could you draw a game tree for this game? If not why not?

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- Simultaneous moves. Our notion of game tree doesn't seem to allow for several players to move simultaneously.
We will see how to deal with these issues later. For now we assume that all our games are of complete (or perfect) information, that is, each player knows precisely where in the game tree they are.


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## Chomp

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The person who has to eat the poisonous square loses. The resulting game is known as $(m \times n)$-Chomp.


## $(2 \times 2)$-Chomp

1 to move

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## $(2 \times 2)$-Chomp



1 to move

2 to move

1 to move

## $(2 \times 2)$-Chomp



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Question. Which player would you rather be, 1 or 2?

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Do you still think this course is going to be interesting?
What would make the games more exciting, and why aren't we looking at those issues?

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The answer to this question depends on whether or not we want to consider symmetry consideration. A crude upper bound for the number of plays ignores these. Player 1 has 9 possibilities for placing his first mark, Player 2 has the remaining 8 squares for her first move, then Player 1 can choose among any of the remaining 7 fields, leading to

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9 \times 8 \times \cdots \times 2 \times 1=9!=362880
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- for each of the branches from such a node there is a label giving the probability of that move occurring.


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To make it simpler, we assume here that the defender has only one army to defend the country with.

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We want to find out whether it is better to attack with one or with two armies.


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1 / 9+1 / 9+1 / 9=3 / 9=1 / 3
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so Attacker wins in $1 / 3$ of all cases.

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Since this is bigger than $1 / 3$, it is better for Attacker to attack with 2 armies. (And, in fact in general for Risk the chance of winning an encounter increases with the number of committed armies.)

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But for this to make sense, the options for the player whose turn it is have to be the same for all the positions he might think he is in.

Question. Why do we make that a requirement?

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In other words, we can cover simultaneous moves by using imperfect information.

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Player 1 has the choice between three moves. Player 2 cannot distinguish between those. Player 2 has three choices for each of these positions, but they are the same in each case. We mark the winner for each possible play.

## Games: Addition to definition

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- for all the nodes in an information set it is the same player's turn, and he is the one who cannot distinguish between them and


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Here is a reminder of our definition of game.

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- a finite set of players,
- a finite game tree,
- for each node of the tree, a player whose turn it is in that position.

As additional information, it is possible to indicate groups of nodes, so called information sets, that one player cannot distinguish between.
The nodes have to have the property that

- for all the nodes in an information set it is the same player's turn, and he is the one who cannot distinguish between them and
- the moves from one of the nodes in the information set are indistinguishable from the moves from any other such node.


## Strategies

In every-day language, people often use the term 'strategy' when they mean 'a general plan to proceed'. For our purposes, we will be looking at a much stricter notion. Assume we are a player in some game.

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Note that nothing in this notion of strategy says that I have to follow some sort of principle! In particular I am allowed to choose different moves to make in positions which look the same on the board, but have different histories.

## Example: Strategy on $(2 \times 2)$-Chomp

Here is an example for a strategy for Player 1 in $(2 \times 2)$-Chomp.

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It is Player 1's turn at the start, so he has to choose a move. It is Player 2's turn next, so Player 1 must allow for all of her possibilities. Again, Player 1 must make a choice for each position he might be in after 2 moves. And finally, Player 2's last potential move has to be accounted for.


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## Example

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## Strategies: towards a definition

Closer inspection shows that a strategy can be described as subtree of the game tree satisfying the following conditions. Assume we have chosen a player for whom we want to give a strategy. Then we can build a subtree, starting from the root, where when we reach a new node,

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- if it is not the chosen player's turn, all available moves are chosen.

Note that we do not worry about positions we can never reach!
Also note that we force a strategy to make a choice of a move if it is the chosen player's turn-we do not allow 'give up by refusing to make a move'. If we want giving up to be an option, we have to make it a move in the game tree.

## Strategies: definition

We can turn this observation into a definition.
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- whenever it is not player $X$ 's turn at a node that belongs to the subtree, all of the available moves (and so all successor nodes) belong to the subtree.
- For all nodes in the same information set for Player $X$, the move chosen by the strategy is the same.


## Different strategies

How many strategies are there for Player 1 in $(2 \times 2)$-Chomp?


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How many strategies are there for Player 1 in $(2 \times 2)$-Chomp? What are they?


1 to move

2 to move

1 to move

2 to move

## Different strategies

How many strategies are there for Player 1 in $(2 \times 2)$-Chomp? What are they?
If we remember that Player 1 starts the game and they then move alternatingly, we can think of the game tree for $(2 \times 2)$-Chomp like this:


## Different strategies

How many strategies are there for Player 1 in $(2 \times 2)$-Chomp? What are they?
Player 1 chooses the left-most move.


## Different strategies

How many strategies are there for Player 1 in $(2 \times 2)$-Chomp? What are they?
Player 1 chooses the second move from the left.


## Different strategies

How many strategies are there for Player 1 in $(2 \times 2)$-Chomp? What are they?
Player 1 chooses the third move from the left.


## Different strategies

How many strategies are there for Player 1 in $(2 \times 2)$-Chomp? What are they?
Player 1 chooses the right-most move. But that means there are further choices down the line! Two decision points with two choices each, that gives $2 \times 2=4$ strategies.


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That makes seven strategies altogether.

## Strategies and imperfect information.

Recall that in our definition of strategy, we stipulated that for all nodes in the same information set for Player $X$, the move chosen by the strategy is the same.

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We consider strategies for Player 2 of Paper-Stone-Scissors.


Player 1

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Player 2

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We consider strategies for Player 2 of Paper-Stone-Scissors. Lastly she may choose to show $S$ for scissors.

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These are all the strategies for Player 2 in Paper-Stone-Scissors.Again they show that Player 2 cannot use information she does not have; her moves for two nodes in the same information set are the same.

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## Strategies and imperfect information.

These are all the strategies for Player 2 in Paper-Stone-Scissors.


So she has the same choices as Player 1, preserving the symmetry of the game as it is defined informally. So we really have not changed the game by using via a game tree.

## Generating all strategies

Let's assume we are given a game tree $t$ and are trying to generate all strategies for Player $X$. This can be done recursively. (Note that when the game is not of perfect information then the following considerations have to be changed accordingly.)

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Note that this does account for chance moves.

## Counting all strategies

Let's assume we are are now trying to count all strategies for Player $X$ for game $t$. This can also be done recursively, with the number of strategies for Player $X$ on game $t$ being denoted by $N_{X}(t)$. (Note that when the game is not of perfect information then once again the following considerations have to be changed accordingly.)

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N_{X}(t)=N_{X}\left(t_{1}\right)+N_{X}\left(t_{2}\right)+\cdots+N_{X}\left(t_{n}\right) .
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## Playing games via strategies

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We then go through the game tree as follows: For every node we reach, starting with the root, we ask the player whose turn it is which move his or her strategy chooses. If there are no chance moves then we thus follow a uniquely determined path through the tree ending at a leaf. That leaf gives us the outcome of playing all these strategies against each other.

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If there are chance moves involved, then we get several paths ending in several leaves, together with a probability for each leaf that it will occur. The probability for a leaf to occur is calculated by multiplying the probabilities occuring along the unique path leading to it. The probabilites for all the possible leaves will add up to one.

## Playing strategies against each other

We have a look at an example to show how playing one strategy against another works.

## Playing strategies against each other

Here is a strategy for $(2 \times 2)$-Chomp for Player 1 .


## Playing strategies against each other

Here is a strategy for $(2 \times 2)$-Chomp for Player 1 . And here is a strategy for the same game for Player 2.


## Playing strategies against each other

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## Playing strategies against each other

Here is a strategy for $(2 \times 2)$-Chomp for Player 1 . And here is a strategy for the same game for Player 2. And here they are together. The outcome of playing the one against the other is immediately obvious.


## Why bother?

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It does have, however, one great disadvantage: It can only be applied if the game is small enough so that one can indeed determine all strategies for all the players.
This is quite hopeless for games like Chess and Go, or even Checkers, simply because the game tree is far too big to hold in a computer-and we've just got some idea of how the number of strategies grows with the size of the game tree!

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This is quite hopeless for games like Chess and Go, or even Checkers, simply because the game tree is far too big to hold in a computer.
Hence when we talk about small games we mean games which are amenable to an analysis via strategies. They will keep us busy for quite a while, maybe longer than you'd thought. This is the most mathematical part of the course.

## Games via strategies

We take the idea of presenting games via their strategies to its logical conclusion.

## Games via strategies

If there are just two players, and no elements of chance involved, then we can present a game in the form of a matrix. We list the strategies for Player 1 along the rows and those for Player 2 along the columns, and record in the centre the outcome of playing the one against the other.

## Games via strategies

In Paper-Stone-Scissors, Player 1 and 2 have three choices each which we can encode as follows.

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In Paper-Stone-Scissors, Player 1 and 2 have three choices each which we can encode as follows. We record the result as

- 1: win for Player 1
- 2: win for Player 2
- D: draw.

|  |  |  | 2 |  |
| :--- | :--- | :--- | :--- | :--- |
|  |  | $P$ | $R$ | $S$ |
|  | $P$ |  |  |  |
| 1 | $R$ |  |  |  |
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|  |  |  |  |  |

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| 1 | $P$ | $D$ | 1 | 2 |
|  | $R$ | 2 | $D$ | 1 |
|  | $S$ | 1 | 2 | $D$ |

Games given likes this are known as matrix games.

## Question

Question. How would you give a matrix version of our miniature Risk game?

## Games via strategies-chance

The problem with matrix games as we've described them so far is that as soon as elements as chance are involved, we don't get a unique result from playing strategies against each other.

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The problem with matrix games as we've described them so far is that as soon as elements as chance are involved, we don't get a unique result from playing strategies against each other.

Recall the miniature Risk game we looked at earlier: Player 1 can only choose between the strategies 'attack with 1' and 'attack with 2', while Player 2 has only one strategy. But there are plenty of possible outcomes for either strategy.


## Games via strategies-chance

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## Games via strategies-chance

The problem with matrix games as we've described them so far is that as soon as elements as chance are involved, we don't get a unique result from playing strategies against each other.

However, if all our outcomes are given as numbers then we can calculate an expected number as the outcome of playing one strategy against another.
If, for each player and each outcome of the game, we have a number to indicate what the outcome means for this player, then we can turn every game into a matrix.

## Pay-off functions

It is for this reason that our definition of 'game' talks about something called a 'pay-off function'.

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Definition 1 A game is given by

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- a finite game tree,
- for each node of the tree, a player whose turn it is in that position and
- for each final node and each player a pay-off function.


## Example

Assume that for Paper-Stone-Scissors the loser pays the winner one (pound, point, chocolate bar), and the pay-off is 0 for both if the game ends in a draw. Then the pay-offs for this game for Player 1 are as follows.

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|  |  |  | 2 |  |
| :---: | :---: | :---: | :---: | :---: |
|  |  | $P$ | $R$ | $S$ |
|  | $P$ | 0 | 1 | -1 |
| 1 | $R$ | -1 | 0 | 1 |
|  | $S$ | 1 | -1 | 0 |

## Example

Assume that for Paper-Stone-Scissors the loser pays the winner one (pound, point, chocolate bar), and the pay-off is 0 for both if the game ends in a draw. Then the pay-offs for this game for Player 1 are as follows. The pay-offs for Player 2 can be given similarly.

|  |  |  | 2 |  |
| :---: | :---: | :---: | :---: | :---: |
|  |  | $P$ | $R$ | $S$ |
|  | $P$ | 0 | 1 | -1 |
| 1 | $R$ | -1 | 0 | 1 |
|  | $S$ | 1 | -1 | 0 |


|  |  |  | 2 |  |
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Question. What would happen if there were more than 2 players?

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|  |  | $P$ | 0 | 1 |


|  |  |  | 2 |  |
| :---: | :---: | :---: | :---: | :---: |
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Note that if we add the entries in the left hand table to the corresponding ones in the right hand table we get a table consisting entirely of 0s. This is because all payments to Player 1 are pay-outs by Player 2 and vice versa.

## Zero-sum games

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Zero-sum games with two players are particularly easy to describe: If we give the pay-off function for one player, we just have to multiply it with -1 to get the pay-off function for the other player.

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Hence in such a case it is sufficient to give just the one matrix, which normally is that for Player 1. In fact, whenever we give a game by just such a matrix, the implicit assumption is that we are describing a 2-person zero-sum game.

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Battle games are another example: Defeating somebody else's forces does not mean that one gets these forces as assets somehow.
Games played in a casino typically aren't zero-sum either-the casino takes a cut which the players never see again.
Some of the most interesting games are non zero-sum. We will see plenty of examples in this course, in particular since most of the known techniques do not work as well for these kinds of games.

## Example

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Consider the following game between two players. Player 1 rolls a a three-faced die. If he throws 1 he pays two units to Player 2. If he throws 2 or 3, Player 2 has a choice. She can either choose to pay one unit to Player 1 (she stops the game) or she can throw the die. If she repeats Player 1's throw, he has to pay her two units. Otherwise she pays him one unit. The game tree is given below, with the pay-off being given for Player 1.

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## Example



Player 1 has only one strategy (he never gets a choice). Player 2 has four strategies (she can choose to throw the die or not, and is allowed to make that dependent on Player 1's throw). We can encode her strategies by saying what she will do when Player 1 throws a 2, and what she will do when Player 1 throws a 3, stop $(S)$ or throw the die $(T)$. So $S \mid T$ means that she will stop if he throws a 2, but throw if he throws a 3. The matrix will look something like this:

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What are the expected pay-offs for the outcome of these strategies? Case $S \mid S$. For each of the possible outcomes of playing this strategy, take the probability that it will occur and multiply it with the pay-off, then add all these up. Hence we get

$$
(1 / 3 \times(-2))+(1 / 3 \times 1)+(1 / 3 \times 1)=0 .
$$

## Example



What are the expected pay-offs for the outcome of these strategies? Case $S \mid T$.
$(1 / 3 \times(-2))+(1 / 3 \times 1)+(1 / 9 \times 1)+(1 / 9 \times 1)+(1 / 9 \times(-2))=-3 / 9=-1 / 3$.

## Example



What are the expected pay-offs for the outcome of these strategies? Case $T \mid S$. This is a symmetric variation of case $S \mid T$, and the pay-off is the same.

## Example



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$$
\begin{array}{r}
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\end{array}
$$

## Example



Question. Which player would you rather be in this game?

A result

We are now ready to state and prove the only result of this introductory section.

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Theorem 1.1 Consider a game with two players, 1 and 2, of perfect information without chance, which can only have three different outcomes: Player 1 wins, Player 2 wins, or they draw. Then one of the following must be true.
(i) Player 1 has a strategy which allows him always to win;
(ii) Player 2 has a strategy which allows him always to win;
(iii) Player 1 and Player 2 both have strategies which ensure that they will not lose (which means that either side can enforce a draw).

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If a player has a strategy which allows him to always win, no matter what the other player does, we call that strategy a winning strategy. Then we can restate the theorem by saying that in such a game, it is the case that either one of the players has a winning strategy or that they can both enforce a draw.

Proof

Induction over the height of the game tree.

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Base case: The height is 0 . Then the result must be stated with the only node, win for Player 1, win for Player 2 or draw, and the Theorem is obviously true.

## Proof

Induction over the height of the game tree.
Induction hypothesis: We can label the root of a game tree of height at most $n$ with a number which indicates which case occurs. We use 1 if Player 1 can force a win, - 1 if Player 2 can force a win and 0 if both sides can enforce a draw.

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Let $t$ be a tree of height $n+1$. By the induction hypothesis we can label the roots of the game trees reached by making a first move of $t$ with a number (say $l_{i}$ for tree $t_{i}$ ) as follows:

- it bears label $l_{i}=1$ if Player 1 wins the game rooted there;
- it bears label $l_{i}=-1$ if Player 2 wins the game rooted there;
- it bears label $l_{i}=0$ if both sides can enforce a draw in the game rooted there.


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There is a tree labelled 1. So there is a $1 \leq i \leq k$ with $l_{i}=1$. Then by making move $m_{i}$ Player 1 can ensure that he will win the subsequent game $t_{i}$ (because he has a winning strategy there), and thus the overall game $t$. Hence $t$ gets label 1 .

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Let $t$ be a tree of height $n+1$. Assume the first move is made by Player 1. The following cases can arise. There is a tree labelled 1. There is no tree labelled 1, but one labelled 0 . So $l_{i} \leq 0$ is true for all $1 \leq i \leq n$ and $l_{i}=0$ for one particular $i$. Then by making the first move $m_{i}$ Player 1 transforms the game into $t_{i}$, where he can ensure a draw.

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There is no tree labelled 1, but one labelled 0 . So $l_{i} \leq 0$ is true for all $1 \leq i \leq n$ and $l_{i}=0$ for one particular $i$. Then by making the first move $m_{i}$ Player 1 transforms the game into $t_{i}$, where he can ensure a draw. But the fact that all the trees have a label of at most 0 means that Player 2 can enforce at least a draw in all games $t_{1}, \ldots, t_{k}$, and thus in the game $t$ no matter what Player 1's first move is. So $t$ gets label 0 .

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Let $t$ be a tree of height $n+1$. Assume the first move is made by Player 1. The following cases can arise.
There is a tree labelled 1.
There is no tree labelled 1, but one labelled 0 .
All trees are labelled -1 . Then no matter which first move $m_{i}$ Player 1 makes, Player 2 can enforce a win in the subsequent game $t_{i}$. Hence she can enforce a win in $t$ and it gets label -1 .

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There is a tree labelled 1.
There is no tree labelled 1, but one labelled 0 .
All trees are labelled -1 .
The case where Player 2 makes the first move is analogous.

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- In 2-player games of perfect information without chance either one of the players can force a win, or they can both force a draw.

