





#### *Janus*: Fine-grained and efficient provenance querying for Taverna

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in collaboration with:

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IMG Seminar, June 2009



- Motivation: Taverna for rapid information integration
- Fine-grained process provenance and its role in Taverna
- Context and scope: forms and uses for process provenance
- Technical challenges in querying provenance traces
- A solution, and experimental results



## Workflow as data integrator



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## The genes $\rightarrow$ pathways workflow in action



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### The genes $\rightarrow$ pathways workflow in action



p2 VEGF signaling pathway

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#### The genes→pathways workflow in action



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#### The genes $\rightarrow$ pathways workflow in action



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goal:

- list all pathways that are mapped from both sets of genes

- substantial list manipulation involved in achieving this

> "p1 MAPK signaling pathway p2 VEGF signaling pathway

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#### The genes→pathways workflow in action



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- list all pathways that are mapped from both sets of genes

- substantial list manipulation involved in achieving this

[[g1, g2, g3], ["g1 p1 p2 ...", "g2 p1 p2 ...", "g3 p1 p2 ..."] [g1, g2, g3]]

> "p1 MAPK signaling pathway p2 VEGF signaling pathway

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## An alternative design





## An alternative design





## An alternative design



## Is either design better?

#### • Pros:

-simpler to design and understand (hopefully)

- (no shims!)
- -accepts multiple gene sets
  - returns list of pathways separately for each gene set
  - in addition to those shared by the union of all sets
- Cons:

-no genes in output list:

[ [ path:mmu04210 Apoptosis, path:mmu04010 MAPK signaling, ...], [ path:mmu04010 MAPK signaling , path:mmu04620 Toll-like receptor, ...] ]

- so the relationship between the gene set and the pathway set is lost...
  - ...Or is it?





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#### Taverna + provenance

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- Taverna type system: strings + nested lists —"cat", ["cat", "dog"], [ ["cat", "dog"], ["large", "small"] ]
- Taverna dataflow model: data-driven execution
   services activate when input is ready
- Workflow provenance: a detailed trace of workflow execution
  - which services were executed
  - when
  - inputs used, outputs produced



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Taverna dataflow model + provenance traces can be a powerful combination

- Causal relations:
  - Which pathway sets come from which gene sets?
  - which processes contributed to producing this image?
  - which process(es) caused this data to be incorrect?
  - which data caused this process to fail?
- Process and data analytics:
  - show me the variations in output in relation to an input parameter sweep (multiple process runs)
  - how often has my favourite service been executed?
    on what inputs?
  - who produced this data?
  - how often does this pathway turn up when the input genes range over a certain set S?

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## Forms of provenance ...

Focus is on the data: the observable outcomes from a process

	raw provenance metadata	provenance metadata + interpretation framework
design	<ul> <li>process structure (workflow graph)</li> <li>history of process composition - reuse</li> <li>process versions</li> </ul>	<ul> <li>service annotations:</li> <li>ex. get_pathways_by_genes</li> <li>who created /edited: attribution</li> <li>why: purpose, intent</li> </ul>
execution	<pre>process events: - service invocation - data production / consumption - causal dependency graphs ex.: - list_of_geneIDList = [ a, b, c] - paths_per_gene = [ [d,e,f], [g,h,j]]  in run #32</pre>	<ul> <li>data annotations,</li> <li>results interpretation in terms of conceptual data model:</li> <li>set of pathways → gene sets</li> </ul>

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#### MANCHESTER 1824 ... and their uses and associated challenges

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$a_1 \cup$
$\sim$ $\sim$
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	raw provenance metadata	provenance metadata + interpretation framework
design	<ul> <li>exploiting semantic properties of the process structure to improve provenance exploitation</li> </ul>	<ul> <li>semantic-based search of process space</li> </ul>
	<ul> <li>exploring process space across versions and structural similarities</li> </ul>	
	<ul> <li>graph matching</li> </ul>	
execution	<ul> <li>enabling partial re-runs of resource-intensive workflows</li> </ul>	<ul> <li>semantic-based query answering over annotated traces</li> </ul>
	<ul> <li>storing very large provenance traces that accumulate over time</li> </ul>	
	- efficient query over large traces	
	- presentation of query answers	



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	- presentation of query answers	The rest of this talk!



#### MANCHESTER 1824 ... and their uses and associated challenges

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## Querying provenance traces

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#### Lineage queries involve traversing a provenance graph from bottom to top





- In most approaches, the originating process are not used for querying
- -consequence: query requires provenance graph traversal
  - large traces  $\rightarrow$  computationally complex
  - view materialization used in practice to get around the computational complexity



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## Requirements for lineage queries - I



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IPAW'08 – Salt Lake City, Utah, June 2008

# Requirements for lineage queries -



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## MANCHESTER Requirements for lineage queries - III

- III Answer queries efficiently without special auxiliary data structures
- (and, please provide declarative query specification)



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Depth mismatch between declared / offered type:

$$depth(P:X2) = depth(c) = 1$$

depth(P:X3) = 1 but depth(c) = 1





Depth mismatch between declared / offered type:

```
depth(P:X1) = 0 but depth(a) = 1
```

```
depth(P:X2) = depth(c) = 1
```

depth(P:X3) = 1 but depth(c) = 1

How *y* is computed at P:

let I = a  $\otimes$  b = [[ <a\_i, b\_j > | b\_j \in b ] | a\_i \in a ] // cross product

I' = [ [ <a\_i, c, b\_j > | b\_j \in b ] |  $a_i \in a$  ] // same product but with c interleaved

$$y = (map (map P) I') = [(map P [  ... ]), ..., (map P [  ... ])] = [[y_{11} ... y_{1n}], ... [y_{n1} ... y_{nm}]]$$





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## Extensional vs intensional tracing

Workflow structure graph



### Extensional vs intensional tracing

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## Extensional vs intensional tracing

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**y**mn

W

bm

b<sub>1</sub>





1) In general the actual depth at the output is:

```
depth(y) = depth(Y) + \sum \delta(X_i = x_i)
```

where  $\delta(X_i = x_i) = depth(x_i) - depth(X_i)$ 

Therefore:

 $\delta(X=x)$  can be computed statically on the workflow graph structure,

- given the declared depth(X)
- using a simple propagation algorithm

#### $Y[i.j] \rightarrow X1[i], X2[], X3[j]$

 $[i_1 . i_2 . ... . i_k] = -------$ 





(0,<mark>2</mark>)

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$$[i_1 . i_2 . ... . i_k] = \frac{\delta(X_1 = x_1)}{k_1 - k_2}$$





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[i<sub>1</sub> . i<sub>2</sub> . . . . i<sub>k</sub>] =

$$\delta(X_2 = x_2)$$

$$X_1$$





(0,<mark>2</mark>)

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X<sub>1</sub> X<sub>2</sub>





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X<sub>1</sub> X<sub>2</sub> X<sub>k</sub>



### Extension to the entire workflow graph



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### Extension to the entire workflow graph

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Y = [ [ ...], ... [...] ]

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### Extension to the entire workflow graph

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 $lineage(P:Y[3,4]) \rightarrow lineage(P:X1[3]),$ lineage(P:X2[]),lineage(P:X2[4]) The University of Manchester MANCHESTER 1824

(0,1)

Х

#### Extension to the entire workflow graph

Q R Y Y (0,1)(1,1) (1,1) **(**0,**1**) (0,1) **X**<sub>3</sub> **X**<sub>1</sub>  $X_2$ [3] Ρ [4] Y

(0,0)

Х

 $lineage(P:Y[3,4]) \rightarrow lineage(P:X1[3]),$ lineage(P:X2[]),lineage(P:X2[4])

Y = [ [ ...], ... [...] ]

(0, 2)

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(0,1)

Х

[3]

## Extension to the entire workflow graph

Ρ

Y

(0, 2)

(0,0)

[4]

Х

Y = [ [ ...], ... [...] ]

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[3]

#### (0,1) (0,0)Х Х [3] Q R Y Y (0,1)(1,1)[3] (1,1) **(**0,**1**) (0,1) **X**<sub>3</sub> **X**<sub>1</sub> $X_2$

Ρ

Y

(0,2)

[4]

Extension to the entire workflow graph

Y = [ [ ...], ... [...] ]

## Extension to the entire workflow graph

[3]

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Y = [ [ ...], ... [...] ]

- lineage(P:Y[3,4])  $\rightarrow$  lineage(P:X1[3]), lineage(P:X2[]), lineage(P:X2[4])
- lineage(P:X1[3]) = lineage(Q:Y[3])  $\rightarrow$ lineage(Q:X[3])

lineage(P:X3[4]) = lineage(R:Y[4])  $\rightarrow$ lineage(R:X[])

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## Extension to the entire workflow graph





## Query processing



- Query processing:
  - alternating sequence
     of *xform* and *xfer* steps
- apply path projection at each *xform* step
- A complete granular and focused query can be answered by traversing the workflow graph alone
  - starting from the target vars
  - one simple query for each selected processor input port



- Scalability:
  - -query time depends on size of workflow graph, not size of provenance graph
  - -workflow graphs are small, fit in memory, can be indexed easily, etc.
  - –search over a graph at least as large as the workflow graph is inevitable -- this is the baseline cost!
- Graceful degradation:
  - -worst case is a completely unfocused query
  - -one query to trace at each *xform* step
  - -no worse than other approaches
- Fine-grain answers provided at no additional cost

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Advantages



## Experimental setup - I

- Performance evaluation performed on programmatically generated dataflows
- -the "T-towers"

#### control:

- size of the lists involved
- length of the paths
- includes one cross product





## Experimental results - I

query response time: naive vs. "path projection" approaches



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## **Experimental results - II**

workflow search time by path length ("tower height") - common to all strategies!



performance degradation on fully unfocused queries 





- An original approach to lineage queries for Taverna that combines
  - efficiency and fine-granularity
- Relies on semantic properties of the Taverna dataflow model
- Further work:
  - visual specification of user query
  - visual presentation of query answer
  - space compression
  - semantic overlays, annotations
- To be bundled with some future version of Taverna...