Implementing DL Systems

Naive Implementations

Problems include:

- Space usage
 - Storage required for tableaux datastructures
 - Rarely a serious problem in practice
- Time usage
 - Search required due to non-deterministic expansion
 - Serious problem in practice
 - Mitigated by:
 - Careful choice of algorithm
 - → Highly optimised implementation

Careful Choice of Algorithm

- Transitive roles instead of transitive closure
 - Deterministic expansion of $\exists R.C$, even when $R \in \mathbf{R}_+$
 - (Relatively) simple blocking conditions
 - Cycles always represent (part of) cyclical models
- Direct algorithm/implementation instead of encodings
 - GCI axioms can be used to "encode" additional operators/axioms
 - Powerful technique, particularly when used with FL closure
 - Can encode cardinality constraints, inverse roles, range/domain, . . .
 - ightharpoonup E.g., (domain R.C) $\equiv \exists R. \top \sqsubseteq C$
 - (FL) encodings introduce (large numbers of) axioms
 - BUT even simple domain encoding is disastrous with large numbers of roles

Highly Optimised Implementation

Optimisation performed at 2 levels

- Computing classification (partial ordering) of concepts
 - Objective is to minimise number of subsumption tests
 - Can use standard order-theoretic techniques
 - → E.g., use enhanced traversal that exploits information from previous tests
 - Also use structural information from KB
 - → E.g., to select order in which to classify concepts
- Computing subsumption between concepts
 - Objective is to minimise cost of single subsumption tests
 - Small number of hard tests can dominate classification time
 - Recent DL research has addressed this problem (with considerable success)

Optimising Subsumption Testing

Optimisation techniques broadly fall into 2 categories

- Pre-processing optimisations
 - Aim is to simplify KB and facilitate subsumption testing
 - Largely algorithm independent
 - Particularly important when KB contains GCI axioms
- Algorithmic optimisations
 - Main aim is to reduce search space due to non-determinism
 - Integral part of implementation
 - But often generally applicable to search based algorithms

Pre-processing Optimisations

Useful techniques include

- Normalisation and simplification of concepts
 - Refinement of technique first used in KRIS system
 - Lexically normalise and simplify all concepts in KB
 - Combine with lazy unfolding in tableaux algorithm
 - Facilitates early detection of inconsistencies (clashes)
- Absorption (simplification) of general axioms
 - Eliminate GCIs by absorbing into "definition" axioms
 - Definition axioms efficiently dealt with by lazy expansion
- Avoidance of potentially costly reasoning whenever possible
 - Normalisation can discover "obvious" (un)satisfiability
 - Structural analysis can discover "obvious" subsumption

Normalisation and Simplification

- Normalise concepts to standard form, e.g.:
 - $\exists R.C \longrightarrow \neg \forall R.\neg C$
 - $C \sqcup D \longrightarrow \neg(\neg C \sqcap \neg D)$
- Simplify concepts, e.g.:
 - $(D \sqcap C) \sqcap (A \sqcap D) \longrightarrow A \sqcap C \sqcap D$
 - $\forall R. \top \longrightarrow \top$
 - $\dots \sqcap C \sqcap \dots \sqcap \neg C \sqcap \dots \longrightarrow \bot$
- Lazily unfold concepts in tableaux algorithm
 - Use names/pointers to refer to complex concepts
 - Only add structure as required by progress of algorithm
 - Detect clashes between lexically equivalent concepts

```
{HappyFather, ¬HappyFather} → clash
{∀has-child.(Doctor ⊔ Lawyer), ∃has-child.(¬Doctor □ ¬Lawyer)} → search
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Absorption I

- Reasoning w.r.t. set of GCI axioms can be very costly
 - GCI $C \sqsubseteq D$ adds $D \sqcup \neg C$ to every node label
 - Expansion of disjunctions leads to search
 - With 10 axioms and 10 nodes search space already 2^{100}
 - GALEN (medical terminology) KB contains hundreds of axioms
- Reasoning w.r.t. "primitive definition" axioms is relatively efficient
 - For $CN \sqsubseteq D$, add D only to node labels containing CN
 - For CN $\supseteq D$, add $\neg D$ only to node labels containing \neg CN
 - Can expand definitions lazily
 - Only add definitions after other local (propositional) expansion
 - Only add definitions one step at a time

Absorption II

- Transform GCIs into primitive definitions, e.g.
 - $\mathsf{CN} \sqcap C \sqsubseteq D \longrightarrow \mathsf{CN} \sqsubseteq D \sqcup \neg C$
 - $CN \sqcup C \supset D \longrightarrow CN \supset D \cap \neg C$
- Absorb into existing primitive definitions, e.g.
 - $\mathsf{CN} \sqsubseteq A$, $\mathsf{CN} \sqsubseteq D \sqcup \neg C \longrightarrow \mathsf{CN} \sqsubseteq A \sqcap (D \sqcup \neg C)$
 - $CN \supseteq A$, $CN \supseteq D \sqcap \neg C \longrightarrow CN \supseteq A \sqcup (D \sqcap \neg C)$
- Use lazy expansion technique with primitive definitions
 - Disjunctions only added to "relevant" node labels
- Performance improvements often too large to measure
 - At least four orders of magnitude with GALEN KB

Algorithmic Optimisations

Useful techniques include

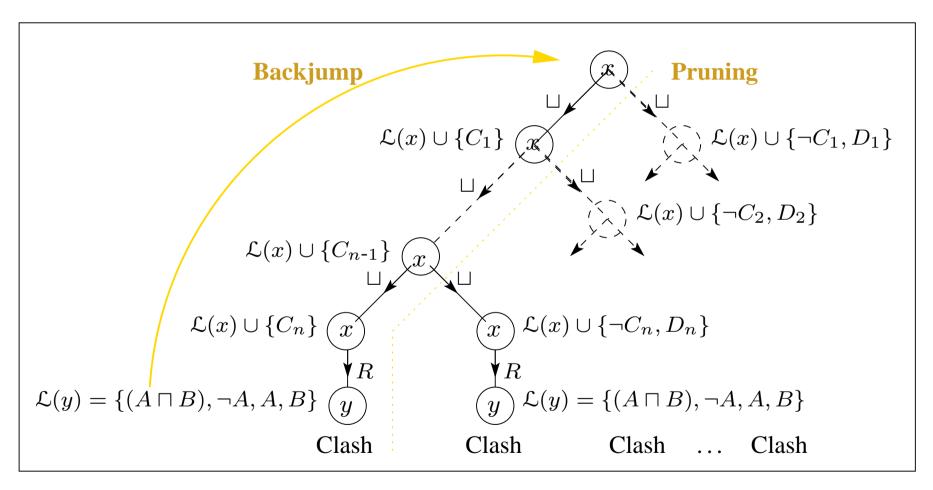
- Avoiding redundancy in search branches
 - Davis-Putnam style semantic branching search
 - Syntactic branching with no-good list
- Dependency directed backtracking
 - Backjumping
 - Dynamic backtracking
- Caching
 - Cache partial models
 - Cache satisfiability status (of labels)
- Heuristic ordering of propositional and modal expansion
 - Min/maximise constrainedness (e.g., MOMS)
 - Maximise backtracking (e.g., oldest first)

Dependency Directed Backtracking

- Allows rapid recovery from bad branching choices
- Most commonly used technique is backjumping
 - Tag concepts introduced at branch points (e.g., when expanding disjunctions)
 - Expansion rules combine and propagate tags
 - On discovering a clash, identify most recently introduced concepts involved
 - Jump back to relevant branch points without exploring alternative branches
 - Effect is to prune away part of the search space
 - Performance improvements with GALEN KB again too large to measure

Backjumping

E.g., if $\exists R. \neg A \sqcap \forall R. (A \sqcap B) \sqcap (C_1 \sqcup D_1) \sqcap \ldots \sqcap (C_n \sqcup D_n) \subseteq \mathcal{L}(x)$



Caching

- Cache the satisfiability status of a node label
 - Identical node labels often recur during expansion
 - Avoid re-solving problems by caching satisfiability status
 - \rightarrow When $\mathcal{L}(x)$ initialised, look in cache
 - Use result, or add status once it has been computed
 - Can use sub/super set caching to deal with similar labels
 - Care required when used with blocking or inverse roles
 - Significant performance gains with some kinds of problem
- Cache (partial) models of concepts
 - Use to detect "obvious" non-subsumption
 - $C \not\sqsubseteq D$ if $C \sqcap \neg D$ is satisfiable
 - ullet $C \sqcap \neg D$ satisfiable if models of C and $\neg D$ can be merged
 - If not, continue with standard subsumption test
 - Can use same technique in sub-problems

Summary

- Naive implementation results in effective non-termination
- Problem is caused by non-deterministic expansion (search)
 - GCIs lead to huge search space
- Solution (partial) is
 - Careful choice of logic/algorithm
 - Avoid encodings
 - Highly optimised implementation
- Most important optimisations are
 - Absorption
 - Dependency directed backtracking (backjumping)
 - Caching
- Performance improvements can be very large
 - E.g., more than four orders of magnitude

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